

EXISTENCE AND NON-EXISTENCE OF ASSOCIATION SCHEMES*

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Abstract. A $(v, n+1)$ -association scheme is a set $S = \{A_0 = I, A_1, \dots, A_n\}$ of $n+1$ symmetric $(0, 1)$ -matrices of order $v \times v$ such that (i) $\sum_{i=0}^n A_i = J$ (the all-ones matrix), (ii) $A_i A_j = \sum_{k=0}^n a_{ijk} A_k$ for $0 \leq i, j \leq n$, where a_{ijk} are non-negative integers. The main purpose of this paper is to study the existence or non-existence of a $(v, n+1)$ -association scheme for certain parameters v and n , especially for the case of v is a power of two.

1. Introduction. An *association scheme with n classes* (or relations) consists of a finite set X of $v \geq 2$ elements together with $n+1$ non-empty relations R_0, R_1, \dots, R_n defined on X which satisfy conditions (R1) to (R4).

(R1) Each R_i is symmetric, i. e. $(x, y) \in R_i$ implies $(y, x) \in R_i$.

(R2) For every $x, y \in X$, $(x, y) \in R_i$ for exactly one i .

(R3) $R_0 = \{(x, x) : x \in X\}$ is the identity relation.

(R4) If $(x, y) \in R_k$, the number of $z \in X$ such that $(x, z) \in R_i$ and $(y, z) \in R_j$ is a constant a_{ijk} depending on i, j, k but not on the particular choice of x and y .

Association schemes were first introduced by statisticians in connection with the design of experiments ([5], [6], [12], [14], [15], [16], [19]), and have since proved very useful in the study of permutation groups ([7], [10], [11], [18]), graphs ([1], [2], [3], [8]), and coding theory ([9], [13]).

In this paper we call an association scheme with n classes on

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a set of v elements a $(v, n + 1)$ -association scheme or simply a $(v, n + 1)$ -scheme. We describe the relations by their adjacency matrices A_i which are the $v \times v$ matrices with rows and columns labeled by the elements of X and defined by

$$(A_i)_{xy} = \begin{cases} 1 & \text{if } (x, y) \in R_i, \\ 0 & \text{otherwise.} \end{cases}$$

The definition of an association scheme is equivalent to saying that the A_i are non-zero $v \times v$ $(0, 1)$ -matrices which satisfy conditions (A1) to (A4). (We often use the set $\mathcal{S} = \{A_0, \dots, A_n\}$ to denote the association scheme.)

$$(A1) \quad \text{Each } A_i \text{ is symmetric, i. e. } A_i = A_i^t.$$

$$(A2) \quad \sum_{i=0}^n A_i = J \text{ (the all-ones matrix).}$$

$$(A3) \quad A_0 = I.$$

$$(A4) \quad A_i A_j = \sum_{k=0}^n a_{ijk} A_k, \quad i, j = 0, 1, \dots, n.$$

(A1), (A2), and (A4) together imply

$$(A5) \quad A_i A_j = A_j A_i, \quad i, j = 0, 1, \dots, n.$$

Summing up the equalities in (A5) for all j , we can see that every A_i has a constant row sums and column sums v_i i. e.

$$(A6) \quad A_i J = J A_i = v_i J, \quad i = 0, 1, \dots, n.$$

Other relations on a_{ijk} are listed below (see [13]).

$$(A7) \quad \begin{aligned} a_{i i 0} &= v_i, \quad a_{i j k} = a_{j i k}, \quad a_{0 j k} = \delta_{j k}, \\ v_k a_{i j k} &= v_i a_{k j i}, \quad \sum_{j=0}^n a_{i j k} = v_i, \\ \sum_{r=0}^n a_{i j r} a_{r k m} &= \sum_{s=0}^n a_{i s m} a_{j k s}. \end{aligned}$$

The existence of special association schemes, such as Hamming schemes, spectral schemes, cyclotomic schemes, Lee schemes, have been extensively studied, see section 2.5 of [9], for some recent results see [20, 21].

The main purpose of this paper is to study conditions on the parameters v and n for which a $(v, n + 1)$ -scheme exists or does not exist. In section 2, we prove some theorems on association schemes. Section 3 discusses constructions of association schemes from association schemes with smaller parameters v and n . Finally,

in section 4, we use these results to study the existence and non-existence of $(v, n + 1)$ -schemes for certain parameters v and n , especially for the case of v is a power of two.

2. Some theorems on association schemes. Suppose $\mathcal{S} = \{A_0, A_1, \dots, A_n\}$ is a $(v, n + 1)$ -scheme, Let \mathcal{S}_i denote the set of all A_j with $v_j = i$. The following two equalities are important in this paper.

$$(2.1) \quad |\mathcal{S}_1| + 2|\mathcal{S}_2| + \dots + v|\mathcal{S}_v| = v.$$

$$(2.2) \quad |\mathcal{S}_1| + |\mathcal{S}_2| + \dots + |\mathcal{S}_v| = n + 1.$$

Note that some \mathcal{S}_i may be empty.

$A_0 = I \in \mathcal{S}_1$. For any $A_i \in \mathcal{S}$, $A_i \in \mathcal{S}_1$ if and only if A_i is a permutation matrix; in this case $A_i = A_i^t = A_i^{-1}$. Moreover, we have the following theorems.

THEOREM 2.1 \mathcal{S}_1 is an abelian group under matrix multiplication. In fact, \mathcal{S}_1 is isomorphic to $(\mathbb{Z}_2)^m$ for some non-negative integer m , and so $|\mathcal{S}_1| = 2^m$.

Proof. For any pair $A_i, A_j \in \mathcal{S}_1$, $A_i A_j = \sum_{k=0}^n a_{ijk} A_k$ implies that $A_i A_j$ has a constant row sums and column sums $\sum_{k=0}^n a_{ijk} v_k$. Since A_i and A_j are permutation matrices, so is $A_i A_j$. These imply that $1 = \sum_{k=0}^n a_{ijk} v_k$, and so all $a_{ijk} = 0$ except $a_{ijk} = v_h = 1$ for exactly one h , i. e. $A_i A_j = A_h \in \mathcal{S}_1$. This proves that \mathcal{S}_1 is closed under matrix multiplication. Also, the identity matrix $I = A_0 \in \mathcal{S}_1$. And for each $A_i \in \mathcal{S}_1$, $A_i^{-1} = A_i \in \mathcal{S}_1$. So \mathcal{S}_1 is a group under matrix multiplication. By (A5), \mathcal{S}_1 is abelian.

The fact that $A_i = A_i^{-1}$ implies that every element of \mathcal{S}_1 is of order 2 except $A_0 = I$. By the Basic Theorem of Abelian Group (see [17], Theorem 4.6), \mathcal{S}_1 is isomorphic to $(\mathbb{Z}_2)^m$ for some non-negative integer m .

THEOREM 2.2 If $A_i \in \mathcal{S}_1$ and $A_j \in \mathcal{S}_h$, then $A_i A_j \in \mathcal{S}_h$.

Proof. We will prove the theorem by induction on h . By Theorem 2.1, the theorem is true for the case of $h = 1$. Suppose the theorem holds for all $h' < h \geq 2$. Similar to the proof of

Theorem 2.1, $A_i A_j = \sum_{k=0}^n a_{ijk} A_k$ implies that $h = v_j = \sum_{k=0}^n a_{ijk} v_k$. Suppose $a_{ijk} = 0$ for all A_k with $v_k \geq h$. Then, by $A_i = A_i^{-1}$,

$$(2.3) \quad A_j = \sum_{v_k < h} a_{ijk} (A_i A_k).$$

By the induction hypothesis, $A_i A_k \in \mathcal{S}_{v_k}$ for all $v_k < h$. (2.3) is impossible since $A_j \in \mathcal{S}_h$ and (A2). Thus $a_{ijk} \geq 1$ for some A_k with $v_k \geq h$. This implies that all $a_{ijk} = 0$ except $a_{ijm} = 1$ and $v_m = h$ for exactly one m , i. e. $A_i A_j = A_m \in \mathcal{S}_h$.

For any \mathcal{S}_h , consider the relation \sim^h on \mathcal{S}_h defined by: $A_j \sim^h A_k$ if and only if $A_i A_j = A_k$ for some $A_i \in \mathcal{S}_1$. Using Theorem 2.1, it is easy to check that \sim^h is an equivalence relation on \mathcal{S}_h . So \mathcal{S}_h is the disjoint union of equivalence classes $[A_j]$. For each $A_j \in \mathcal{S}_h$, denote $F(A_j) = \{A_i \in \mathcal{S}_1 : A_i A_j = A_j\}$.

THEOREM 2.3 *Suppose $A_j \in \mathcal{S}_h$ and $|\mathcal{S}_1| = 2^m$ as shown in Theorem 2.1. The following statements hold.*

- (i) $F(A_j)$ is a subgroup of \mathcal{S}_1 , and so $F(A_j) \simeq (Z_2)^{m'}$ for some $m' \leq m$.
- (ii) $|[A_j]| = 2^{m-m'}$.
- (iii) $|F(A_j)| \leq h$.
- (iv) $|[A_j]|$ is a multiple of $2^{m-m''}$, where $m'' = \lfloor \log_2 h \rfloor$.
- (v) $|\mathcal{S}_h| = a 2^{m-m''}$ and $h|\mathcal{S}_h| \geq a 2^m$ for some positive integer a .

Proof. (i) Suppose $A_i, A_k \in F(A_j)$, i. e. $A_i, A_k \in \mathcal{S}_1$ and $A_i A_j = A_k A_j = A_j$. Since $A_i^{-1} = A_i \in \mathcal{S}_1$ and $A_i^{-1} A_j = A_i A_j = A_j$, $A_i^{-1} \in F(A_j)$. $A_i A_k \in \mathcal{S}_1$ by Theorem 2.1 and $(A_i A_k) A_j = A_i (A_k A_j) = A_i A_j = A_j$, so $A_i A_k \in F(A_j)$. Also $I = A_0 \in F(A_j)$. These prove that $F(A_j)$ is a subgroup of \mathcal{S}_1 . $\mathcal{S}_1 = (Z_2)^m$ implies $\mathcal{S}(A_j) \simeq (Z_2)^{m'}$ for some $m' \leq m$.

(ii) Note that $[A_j] = \{A_i A_j : A_i \in \mathcal{S}_1\}$. $A_i A_j = A_k A_j$ if and only if $(A_k A_i) A_j = A_j$ or $A_k A_i \in F(A_j)$. So $|[A_j]| = |\mathcal{S}_1|/|F(A_j)| = 2^{m-m'}$.

(iii) By (A6) and (A4),

$$\begin{aligned} v_j J &= J A_j = \sum_{i=0}^n A_i A_j = \sum_{A_i \in F(A_j)} A_i A_j + \sum_{A_i \notin F(A_j)} A_i A_j \\ &= |F(A_j)| A_j + N, \end{aligned}$$

where N is a non-negative matrix. Compare the entries which are ones in A_j , then $h = v_j \geq |F(A_j)|$.

(iv) By (ii) and (iii), $2^{m'} = |F(A_j)| \leq h$, so $m' \leq \log_2 h$. Then $m' \leq m''$ and $|[A_j]| = 2^{m-m'}$ is a multiple of $2^{m-m''}$.

(v) $|\mathcal{S}_h| = a 2^{m-m''}$ follows from (iv) and the fact that \mathcal{S}_h is the disjoint union of equivalence classes $[A_j]$. $h \geq 2^{m''}$ implies $h|\mathcal{S}_h| \geq a 2^m$.

THEOREM 2.4 *Suppose $E \subseteq \{1, \dots, v\}$. If*

$$\sum_{j \in E} |\mathcal{S}_j| = p = 2^{u_s} + \dots + 2^{u_1},$$

where $u_s > u_{s-1} > \dots > u_1 \geq 0$ are integers, then $\sum_{j \in E} j|\mathcal{S}_j| \geq s 2^m$, where $2^m = |\mathcal{S}_1|$.

Proof. We will prove the theorem by induction on p . Suppose $p = 1$, then $p = 2^0$ and there is some $h \in E$ such that $\mathcal{S}_h = \emptyset$. By Theorem 2.3 (v), $h|\mathcal{S}_h| \geq 2^m$. So the theorem holds.

Suppose the theorem holds for all $p' < p = 2^{u_s} + \dots + 2^{u_1} \geq 2$. Since $u_1 \geq m$ would imply $\sum_{j \in E} j|\mathcal{S}_j| \geq \sum_{j \in E} |\mathcal{S}_j| \geq s 2^m$, without loss of generality we assume that $u_1 < m$. Suppose $j < 2^{m-u_1}$ for all $j \in E$. By Theorem 2.3 (v), each $|\mathcal{S}_j|$ is a multiple of $2^{m-m''}$, where $m'' = \lfloor \log_2 j \rfloor < m - u_1$, i. e. $m - m'' \geq u_1 + 1$. Thus $p = \sum_{j \in E} |\mathcal{S}_j|$ is a multiple of 2^{u_1+1} which contradicts $p = 2^{u_s} + \dots + 2^{u_2} + 2^{u_1}$ and $u_s > \dots > u_2 \geq u_1 + 1 > u_1$. So there is some $h \in E$ with $h \geq 2^{m-u_1}$. By Theorem 2.3(v), $|\mathcal{S}_h| = a 2^{m-m''}$ and $h|\mathcal{S}_h| \geq a 2^m$, where a is a positive integer and $m'' = \lfloor \log_2 h \rfloor \geq m - u_1$, i. e. $2^{m-m''} \leq 2^{u_1}$.

Consider $E' = E \setminus \{h\}$, then $p' = p - |\mathcal{S}_h|$, where $|\mathcal{S}_h| \leq a 2^{u_1} \leq 2^{u_s} + 2^{u_{s-1}} + \dots + 2^{u_1}$. So $p' = 2^{u_s} + \dots + 2^{u_{a+1}} + 2^{w_r} + \dots + 2^{w_1}$ where $r \geq 0$ and $u_s > \dots > u_{a+1} > w_r > \dots > w_1 \geq 0$. By the induction hypothesis, $\sum_{j \in E'} j|\mathcal{S}_j| \geq (s - a + r) 2^m \geq (s - a) 2^m$. So $\sum_{j \in E} j|\mathcal{S}_j| \geq a 2^m + (s - a) 2^m = s 2^m$.

Hence the theorem holds by induction.

THEOREM 2.5 $|\mathcal{S}_1|$ is a divisor of v .

Proof. Consider the binary relation \sim on X defined by: $i \sim j$ if and only if there is a matrix $A_k = (a_{ij}^{(k)}) \in \mathcal{S}_1$ such that $a_{ij}^{(k)} = 1$. \sim is an equivalent relation as shown below.

(i) $i \sim i$ since $A_0 = I = (\delta_{ij})$ with $\delta_{ii} = 1$.

(ii) $i \sim j$ implies $j \sim i$ follows from the fact that each A_k is symmetric.

(iii) Suppose $i \sim j$ and $j \sim k$, i. e. there are $A_p, A_q \in \mathcal{S}_1$ with $a_{ij}^{(p)} = a_{jk}^{(q)} = 1$. Note that $A_r = A_p A_q \in \mathcal{S}_1$ and $a_{ik}^{(r)} \geq a_{ij}^{(p)} a_{jk}^{(q)} = 1$. So $a_{ik}^{(r)} = 1$ and then $i \sim k$.

The equivalence relation \sim partitions X into equivalence classes. For each $i \in X$ and $A_k \in \mathcal{S}_1$, there exists one $j \in [i]$ such that $a_{ij}^{(k)} = 1$. The correspondence between A_k and j is one to one. Thus each class $[i]$ is of size $|\mathcal{S}_1|$. Then $v = |\mathcal{S}_1| \cdot \#(\text{equivalence classes})$, i. e. $|\mathcal{S}_1|$ is a divisor of v .

3. Constructions of new association schemes from given association schemes. In coding theory, *extension* of an association scheme is a very common way to get new association schemes (see [9], section 2.5). In this section, we introduce other constructions of new association schemes from given association schemes by means of Kronecker product of two matrices.

Suppose $A = (a_{ij})$ is a $u \times v$ matrix and B a $r \times s$ matrix. The *Kronecker product* $A \otimes B$ of A and B is the following $(ur) \times (vs)$ matrix

$$A \otimes B = \begin{pmatrix} a_{11} B & a_{12} B & \cdots & a_{1v} B \\ a_{21} B & a_{22} B & \cdots & a_{2v} B \\ \vdots & \vdots & \ddots & \vdots \\ a_{u1} B & a_{u2} B & \cdots & a_{uv} B \end{pmatrix}.$$

The following equalities are frequently used in this paper.

$$(3.1) \quad a(A \otimes B) = (aA) \otimes B = A \otimes (aB).$$

$$(3.2) \quad (A_1 + A_2) \otimes B = (A_1 \otimes B) + (A_2 \otimes B).$$

$$(3.3) \quad A \otimes (B_1 + B_2) = (A \otimes B_1) + (A \otimes B_2).$$

$$(3.4) \quad (A \otimes B)(C \otimes D) = (AC) \otimes (BD).$$

$$(3.5) \quad (A \otimes B)^t = A^t \otimes B^t.$$

THEOREM 3.1 *Suppose $\mathcal{S} = \{A_0, \dots, A_n\}$ is a $(v, n + 1)$ -scheme $\mathcal{S}' = \{A'_0, \dots, A'_m\}$ is a set of $m + 1$ $v \times v$ $(0, 1)$ -matrices such that (i) each A'_j is the sum of some matrices in \mathcal{S} , (ii) all A'_j sum up to J , (iii) $A_i A'_j = \sum_{k=0}^m a'_{ijk} A'_k$ for $0 \leq i \leq n$ and $0 \leq j \leq m$, where a'_{ijk} are non-negative integers. $\mathcal{G} = \{B_0, \dots, B_p\}$ is a $(u, p + 1)$ -scheme. $\mathcal{G}' \subseteq \mathcal{G}$ satisfies $B_0 \in \mathcal{G}'$ and $B_i B_j = \sum_{B_k \in \mathcal{G}'} b'_{ijk} B_k$ for $B_i, B_j \in \mathcal{G}'$, where b'_{ijk} are non-negative integers. Construct the set*

$$\mathcal{U} = \{A_i \otimes B_j : A_i \in \mathcal{S} \text{ and } B_j \in \mathcal{G}'\} \cup \{A'_i \otimes B_j : A'_i \in \mathcal{S}' \text{ and } B_j \in \mathcal{G} \setminus \mathcal{G}'\}.$$

Then \mathcal{U} is an association scheme.

Proof. We shall prove that (A1) to (A4) hold for \mathcal{U} .

(A1) Since all matrices in $\mathcal{S}, \mathcal{S}', \mathcal{G}, \mathcal{G}'$ are symmetric nonzero $(0, 1)$ -matrices, so are all matrices in \mathcal{U} .

$$\begin{aligned} \text{(A2)} \quad & \sum \{A_i \otimes B_j : A_i \in \mathcal{S} \text{ and } B_j \in \mathcal{G}'\} \\ & + \sum \{A'_i \otimes B_j : A'_i \in \mathcal{S}' \text{ and } B_j \in \mathcal{G} \setminus \mathcal{G}'\} \\ & = \sum_{A_i \in \mathcal{S}} A_i \otimes \sum_{B_j \in \mathcal{G}'} B_j + \sum_{A'_i \in \mathcal{S}'} A'_i \otimes \sum_{B_j \in \mathcal{G} \setminus \mathcal{G}'} B_j \\ & = J_v \otimes \sum_{B_j \in \mathcal{G}'} B_j + J_v \otimes \sum_{B_j \in \mathcal{G} \setminus \mathcal{G}'} B_j \\ & = J_v \otimes \left(\sum_{B_j \in \mathcal{G}'} B_j + \sum_{B_j \in \mathcal{G} \setminus \mathcal{G}'} B_j \right) \\ & = J_v \otimes J_u \\ & = J_{vu}. \end{aligned}$$

$$\text{(A3)} \quad I_{vu} = I_v \otimes I_u = A_0 \otimes B_0 \in \mathcal{U}.$$

$$\text{(A4)} \quad \text{For } A_i \in \mathcal{S}, B_j \in \mathcal{G}', A_r \in \mathcal{S}, B_s \in \mathcal{G}',$$

$$\begin{aligned} & (A_i \otimes B_j)(A_r \otimes B_s) \\ & = (A_i A_r) \otimes (B_j B_s) \\ & = \sum_{A_k \in \mathcal{S}} a_{irk} A_k \otimes \sum_{B_h \in \mathcal{G}'} b'_{jsh} B_h \\ & = \sum \{a_{irk} b'_{jsh} A_k \otimes B_h : A_k \in \mathcal{S} \text{ and } B_h \in \mathcal{G}'\}. \end{aligned}$$

For $A_i \in \mathcal{S}$, $B_j \in \mathcal{G}'$, $A_r \in \mathcal{S}'$, $B_s \in \mathcal{G} \setminus \mathcal{G}'$,

$$\begin{aligned} & (A_i \otimes B_j)(A_r \otimes B_s) \\ &= (A_i A_r) \otimes (B_j B_s) \\ &= \sum_{A'_k \in \mathcal{S}'} a'_{irk} A'_k \otimes \sum_{B_h \in \mathcal{G}} b_{jsh} B_h \\ &= \sum \{a'_{irk} b_{jsh} A'_k \otimes B_h : A'_k \in \mathcal{S}' \text{ and } B_h \in \mathcal{G} \setminus \mathcal{G}'\} \\ &\quad + \sum \{a'_{irk} b_{jsh} c_{kt} A_t \otimes B_h : A_t \in \mathcal{S} \text{ and } B_h \in \mathcal{G}'\} \end{aligned}$$

where $A'_k = \sum_{t=0}^n c_{kt} A_t$ by assumption (i) of the theorem.

For $A_i \in \mathcal{S}'$, $B_j \in \mathcal{G} \setminus \mathcal{G}'$, $A_r \in \mathcal{S}'$, $B_s \in \mathcal{G} \setminus \mathcal{G}'$, we can similarly prove that $(A_i \otimes B_j)(A_r \otimes B_s)$ is a non-negative integer combination of matrices in \mathcal{U} .

COROLLARY 3.2 *Suppose $\mathcal{S} = \{A_0, \dots, A_n\}$ is a $(v, n + 1)$ -scheme and $\mathcal{G} = \{B_0, \dots, B_p\}$ is a $(u, p + 1)$ -scheme. Then $\mathcal{S} \otimes \mathcal{G} = \{A_i \otimes B_j : 0 \leq i \leq n \text{ and } 0 \leq j \leq p\}$ is a $(vu, (n + 1)(p + 1))$ -scheme.*

Proof. Choose $\mathcal{S}' = \mathcal{S}$ and $\mathcal{G}' = \{B_0\}$. Apply Theorem 3.1.

$\mathcal{S} \otimes \mathcal{G}$ is called the *type I product* of \mathcal{S} and \mathcal{G} . Note that Hamming scheme of length 3 is in fact equivalent to $\{I_4, J_4 - I_4\} \otimes \{I_2, J_2 - I_2\}$.

COROLLARY 3.3 *Suppose $\mathcal{S} = \{A_0, \dots, A_n\}$ is a $(v, n + 1)$ -scheme and $\mathcal{G} = \{B_0, \dots, B_p\}$ is a $(u, p + 1)$ -scheme. $\mathcal{S}' = \{J_v\}$ and \mathcal{G}' is a subgroup of $\mathcal{G}_1 = \{B_j : \text{row sum of } B_j \text{ is } 1\}$. Then $\{A_i \otimes B_j : A_i \in \mathcal{S} \text{ and } B_j \in \mathcal{G}'\} \cup \{J_v \otimes B_j : B_j \in \mathcal{G} \setminus \mathcal{G}'\}$ is a $(vn, n|\mathcal{G}'| + p + 1)$ -scheme. (This new scheme is called the *type II product* of \mathcal{S} and \mathcal{G} with respect to \mathcal{G}' .)*

4. Existence or non-existence of association schemes with certain parameters v and n . The main purpose of this section is to prove the existence or non-existence of schemes with certain parameters v and n . In other word, we want to determine N_v , the set of integers $n + 1$ for which there exists a $(v, n + 1)$ -scheme. It is easy to see that $2 \leq \min N_v \leq \max N_v \leq v$. The only $(v, 2)$ -scheme is $\{I, J - I\}$. so $\min N_v = 2$.

A $(v, 3)$ -scheme $\{A_0, A_1, A_2\}$ is equivalent to a *strongly regular* graph whose adjacency matrix is A_1 (see [4]). For v a composite

integer or a prime of $4s + 1$ type, there always exists a $(v, 3)$ -scheme. But it is unsolved for the case of v is a prime of $4s + 3$ type. We know that there is no $(3, 3)$ -scheme or $(7, 3)$ -scheme.

If $v = rs$, with $r, s \geq 2$ are integers, then $\{I_v, (J_r - I_r) \otimes I_s, J_r \otimes (J_s - I_s)\}$ is a $(v, 3)$ -scheme. In fact this scheme is the type II product of $\{I_r, J_r - I_r\}$ and $\{I_s, J_s - I_s\}$ with respect to $\{I_s\}$. If v is a prime of $4s + 1$ type, we can use the fact that Z_v has $2s$ quadratic residue to construct a (v, s) -scheme.

Next question is to determine $\max N_v$ in term of v .

THEOREM 4.1 *If $v = 2^r u$, where r is a non-negative integer and u an odd integer, then $\max N_v = 2^{r-1}(u + 1)$.*

Proof. Suppose \mathcal{S} is a $(v, n + 1)$ -scheme. By Theorems 2.1 and 2.5, $|\mathcal{S}_1| = 2^m$ with $m \leq r$. By (2.1) and (2.2),

$$v \geq |\mathcal{S}_1| + 2 \sum_{j=2}^v |\mathcal{S}_j| = 2^m + 2(n + 1 - 2^m) = 2(n + 1) - 2^m.$$

Then

$$n + 1 \leq (2^r u + 2^m)/2 \leq (2^r u + 2^r)/2 = 2^{r-1}(u + 1).$$

So $\max N_v \leq 2^{r-1}(u + 1)$.

Conversely, we will construct a $(v, 2^{r-1}(u + 1))$ -scheme and conclude that $\max N_v = 2^{r-1}(u + 1)$.

Consider the following $u \times u$ permutation matrices $P_k = (t_{ij}^{(k)})$, $0 \leq k \leq u - 1$, defined by $t_{ij}^{(k)} = 1$ if $j = i + k$ and $t_{ij}^{(k)} = 0$ otherwise, where the addition of indices are taken modulo u .

let $A_0 = P_0$ and $A_i = P_i + P_{u-i}$ for $1 \leq i \leq (u - 1)/2$. It is straight forward to check that $\mathcal{S} = \{A_0, A_1, \dots, A_{(u-1)/2}\}$ is a $(u, (u + 1)/2)$ -scheme by using the fact that $P_k P_h = P_{k+h}$. Consider the $(2, 2)$ -scheme $\mathcal{G} = \{I_2, J_2 - I_2\}$. Then $\mathcal{S} \otimes \mathcal{G} \otimes \dots \otimes \mathcal{G}$ (with r terms of \mathcal{G}) is a $(v, 2^{r-1}(u + 1))$ -scheme.

In the rest of this section, we will concentrate on the case of $v = 2^n$.

THEOREM 4.2 *Suppose $v = 2^n$ and $n + 1 = 2^{u_0} + 2^{u_1} + \dots + 2^{u_r} \geq 2$, where $u_0 > u_1 > \dots > u_r \geq 0$ are integers. If $u \geq r + u_0$, then there exists a $(v, n + 1)$ -scheme.*

Proof. We will prove the theorem by induction on u . The case of $u = 1$ is clear since $v = n + 1 = 2$. Suppose the theorem holds for all $u' < u \geq 2$. Without loss of generality we can assume that $n + 1 \geq 3$, since $\{I_v, J_v - I_v\}$ is a $(v, 2)$ -scheme.

Suppose $u_r \geq 1$, i.e. $n + 1$ is even. Consider $v' = 2^{u-1}$ and $n' + 1 = (n + 1)/2$. Note that $n' + 1 = 2^{u_0-1} + 2^{u_1-1} + \dots + 2^{u_r-1} \geq 2$. Since $u > r + u_0$ implies $u - 1 \geq r + (u_0 - 1)$, by the induction hypothesis, there is a $(v', n' + 1)$ -scheme \mathcal{S} . Consider the $(2, 2)$ -scheme $\mathcal{G} = \{I_2, J_2 - I_2\}$. By corollary 3.2, $\mathcal{S} \otimes \mathcal{G}$ is a $(v, n + 1)$ -scheme.

Suppose $u_r = 0$, i.e. $n + 1$ is odd. for the case of $n + 1 = 3$, a $(v, 3)$ -scheme exists as shown in the 3rd paragraph of this section. For the case of $n + 1 \geq 5$, consider $v' = 2^{u-2}$ and $n' + 1 = n/2$. Note that $n' + 1 = 2^{u_0-1} + 2^{u_1-1} + \dots + 2^{u_{r-1}-1} \geq 2$. Since $u > r + u_0$ implies $u - 2 \geq (r - 1) + (u_0 - 1)$, by the induction hypothesis, there is a $(v', n' + 1)$ -scheme \mathcal{S} . Next consider the $(4, 3)$ -scheme $\mathcal{G} = \{I_4, (J_2 - I_2) \times I_2, J_2 \otimes (J_2 - I_2)\}$. Let $\mathcal{G}' = \{I_4, (J_2 - I_2) \otimes I_2\}$. By Corollary 3.3, the type II product of \mathcal{S} and \mathcal{G} with respect to \mathcal{G}' is a $(v, n + 1)$ -scheme.

Thus the theorem holds by induction.

Although Theorem 4.2 is proved by induction, we can in fact construct the corresponding $(v, n + 1)$ -scheme in the proof. For convenience, we use the following notation.

$$I = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, \quad J = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}, \quad K = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix},$$

$$\langle M_1, \dots, M_u \rangle = M_1 \otimes \dots \otimes M_u.$$

Suppose $v = 2^u$ and $n + 1 = 2^{u_0} + 2^{u_1} + \dots + 2^{u_r} \geq 2$, where $u_0 > u_1 > \dots > u_r = 0$ are integers and $u \geq r + u_0$. (If $u_r > 0$, we can consider $v' = 2^{u-u_r}$ and $n' + 1 = (n + 1)/2^{u_r}$. Then use type I product of schemes.) Let $u_i = 0$ for $i > r$. Define

$$\mathcal{S}^{(0)} = \{ \langle M_1, \dots, M_u \rangle :$$

$$M_j = I \text{ or } K \text{ for } 1 \leq j \leq u_0, M_j = I \text{ for } j \geq u_0 + 1 \},$$

$$\mathcal{S}^{(i)} = \{ \langle M_1, \dots, M_u \rangle : M_j = J \text{ for } 1 \leq j \leq u_0 + i - u_i - 1,$$

$$M_j = I \text{ or } K \text{ for } u_0 + i - u_i \leq j \leq u_0 + i - 1,$$

$$M_{u_0+i} = K, M_j = I \text{ for } j \geq u_0 + i + 1 \},$$

$$\text{for } i = 1, 2, \dots, u - u_0.$$

Then $\mathcal{S} = (\cup_{0 \leq i \leq r-1} \mathcal{S}^{(i)}) \cup \{A_n\}$ is the scheme given in the proof of Theorem 4.2, where $A_n = \sum \{M : \mathcal{S}^{(i)} = \{M\}, r \leq i \leq u - u_0\}$.

For example, $v = 2^8$ and $n + 1 = 2^4 + 2^2 + 2^1 + 2^0$. Then \mathcal{S} contains the following $n + 1$ matrices.

$$\begin{aligned} &\langle K^{c_1}, K^{c_2}, K^{c_3}, K^{c_4}, I, I, I, I \rangle, & c_1, c_2, c_3, c_4 &= 1 \text{ or } 2, \\ &\langle J, J, K^{c_5}, K^{c_6}, K, I, I, I \rangle, & c_5, c_6 &= 1 \text{ or } 2, \\ &\langle J, J, J, J, K^{c_7}, K, I, I \rangle, & c_7 &= 1 \text{ or } 2, \\ &\langle J, J, J, J, J, J, K, I \rangle + \langle J, J, J, J, J, J, J, K \rangle, \end{aligned}$$

where $K^1 = K$ and $K^2 = I$.

COROLLARY 4.3 *Suppose $v = 2^u$ and $2 \leq n + 1 \leq 2^{1+\lceil u/2 \rceil}$, then there exists a $(v, n + 1)$ -scheme.*

THEOREM 4.4 *Suppose \mathcal{S} is a $(v, n + 1)$ -scheme with $2^{u-1} < n + 1 < v = 2^u$, then $n + 1 = 2^{u-1} + 2^w$ for some $0 \leq w \leq u - 2$.*

Proof. Suppose the theorem is not true, then $n + 1 = 2^{u-1} + 2^w + z$, where $0 \leq w \leq u - 2$ and $1 \leq z \leq 2^w - 1$. By (2.1) and (2.2),

$$\begin{aligned} v &\geq |\mathcal{S}_1| + 2 \sum_{j=2}^v |\mathcal{S}_j| \\ &= |\mathcal{S}_1| + 2(n + 1 - |\mathcal{S}_1|) = 2(n + 1) - |\mathcal{S}_1|. \end{aligned}$$

Then

$$|\mathcal{S}_1| \geq 2(n + 1) - v = 2^{w+1} + 2z > 2^{w+1}.$$

By Theorem 2.1, $|\mathcal{S}_1| = 2^m$ for some non-negative integer m . $2^u = v > n + 1 \geq 2^m > 2^{w+1}$ implies $u - 1 \geq m \geq w + 2$. Again, by (2.1) and (2.2),

$$\begin{aligned} 2^u = v &\geq |\mathcal{S}_1| + 2|\mathcal{S}_2| + 3 \sum_{j=3}^v |\mathcal{S}_j| \\ &= |\mathcal{S}_1| + 2|\mathcal{S}_2| + 3(n + 1 - |\mathcal{S}_1| - |\mathcal{S}_2|) \\ &= 3 \cdot 2^{u-1} + 3 \cdot 2^w + 3z - 2^{m+1} - |\mathcal{S}_2|. \end{aligned}$$

Then

$$|\mathcal{S}_2| \geq 2^{u-1} - 2^{m+1} + 3 \cdot 2^w + 3z.$$

By Theorem 2.3(v), $|\mathcal{S}_2|$ is a multiple of 2^{m-1} . Therefore, $|\mathcal{S}_2| = 2^{u-1} - 2^{m+1} + t \cdot 2^{m-1}$, where $t \geq 1$ is an integer. So

$$(4.1) \quad \sum_{j=3}^v |\mathcal{S}_j| = n + 1 - |\mathcal{S}_1| - |\mathcal{S}_2| \\ = 2^w + z + (2 - t) 2^{m-1} \geq 0$$

and

$$(4.2) \quad \sum_{j=3}^v j |\mathcal{S}_j| = v - |\mathcal{S}_1| - 2|\mathcal{S}_2| = (3 - t)2^m.$$

(4.1) and the fact that $m \geq w + 2$ imply $2 - t \geq 0$, i. e. $2 - t = 0$ or 1. Then

$$\sum_{j=3}^v |\mathcal{S}_j| = (2 - t) 2^{m-1} + 2^w + 2^{u_s} + \dots + 2^{u_1},$$

where $m - 1 > w > u_s > \dots > u_1 \geq 0$ and $s \geq 1$. (Note that $(2 - t) 2^{m-1}$ is either 2^{m-1} or nothing.) By Theorem 2.4,

$$\sum_{j=3}^v j |\mathcal{S}_j| \geq (2 - t + 1 + s)2^m \geq (4 - t)2^m.$$

This contradicts (4.2). Thus the theorem holds.

By Corollary 4.3 and Theorem 4.4, we have

$$N_2 = \{2\}, \\ N_4 = \{2, 3, 4\}, \\ N_8 = \{2, 3, 4, 5, 6, 8\}, \\ N_{16} = \{2, 3, 4, 5, 6, 7, 8, 9, 10, 12, 16\}.$$

However, for $v = 32$, the only $n + 1$ for which we can not use Theorems 4.2 and 4.4 to determine if $n + 1 \in N_v$ is $n + 1 = 15$. But there exists (4, 3)-scheme and (8, 5)-scheme. Their type I product is a (32, 15)-scheme. So we have

$$N_{32} = \{2, 3, \dots, 16, 17, 18, 24, 32\}.$$

Similary, for $v = 2^6$, we have difficulty when $n + 1 = 23, 27, 29, 30, 31$. Since there exist (4, 3)-scheme, (16, 9)-scheme, (8, 5)-scheme, (8.6)-scheme, by using type I product, there exist (32, 27)-scheme and (32, 30)-scheme. Also the following is a (32, 23)-scheme:

$$\mathcal{S} = \{ \langle K^{c_1}, K^{c_2}, K^{c_3}, K^{c_4}, I, I \rangle, \quad c_1, c_2, c_3, c_4 = 1 \text{ or } 2, \\ \langle J, J, K^{c_5}, K^{c_6}, K, I \rangle, \quad c_5, c_6 = 1 \text{ or } 2, \\ \langle K^{c_7}, J, J, J, I, K \rangle, \quad c_7 = 1 \text{ or } 2, \\ \langle J, J, J, J, K, K \rangle \}.$$

so only the cases of $n + 1 = 29, 31$ are unknown.

$$N_{64} = \{2 \cdots 28, (29?), 30, (31?) 32, 33, 34, 36, 40, 48, 64\}.$$

In general, we still can not determine N_v . In order to get more results on N_v , we believe that we should understand more structures on \mathcal{S}_j for an association scheme \mathcal{S} .

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